

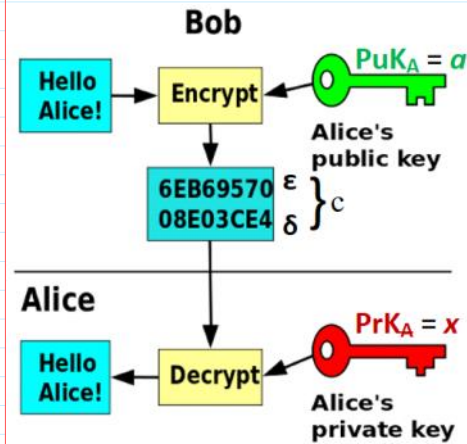
Course for repetition:

<http://crypto.fmf.ktu.lt/telekonf/archyvas/B111%20Kriptologija/B111%202024-R/>

### Asymmetric Encryption - Decryption

$$c = \text{Enc}(\text{PuK}_A, m)$$

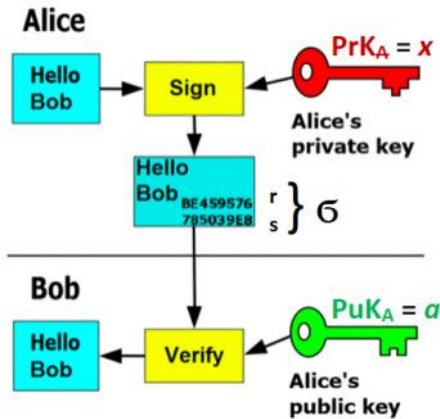
$$m = \text{Dec}(\text{PrK}_A, c)$$



### Asymmetric Signing - Verification

$$\text{Sign}(\text{PrK}_A, m) = \sigma = (r, s)$$

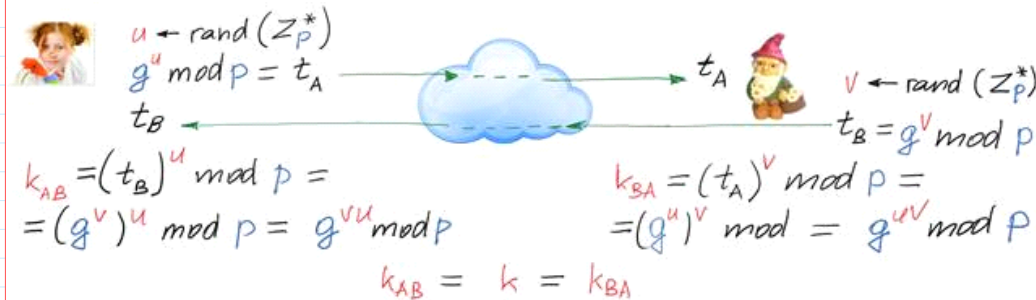
$$V = \text{Ver}(\text{PuK}_A, m, s), V \in \{\text{True}, \text{False}\} \equiv \{1, 0\}$$



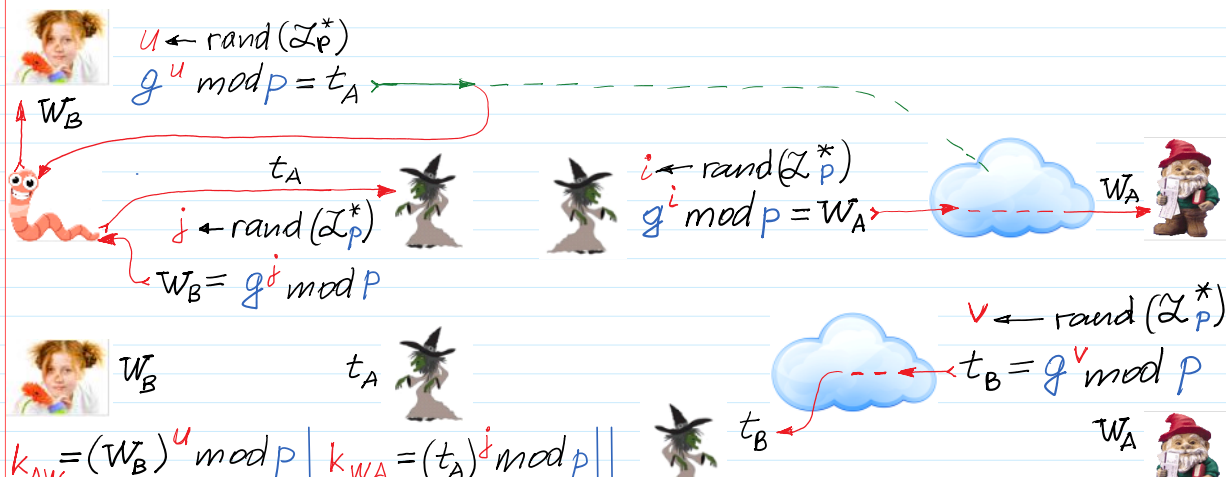
> Documents > 100 MOKYMAS  
 1 DLF v-4.pptx  
 2 MIMA.pptx

### Diffie-Hellman Key Agreement Protocol (DH KAP)

Public Parameters  $PP = (p, g)$ ;  $Z_p^* = \{1, 2, 3, \dots, p-1\}$ ;  $*$  mod  $p$ .

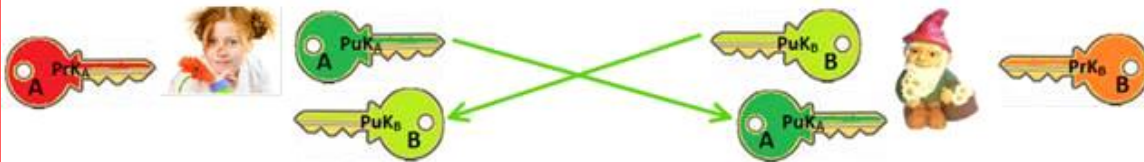


### Man in the Middle Attack - MiMA - Impersonation Attack



$$\begin{array}{l}
 k_{AW} = (W_B)^u \bmod p \\
 = (g^j)^u \bmod p \\
 = g^{ju} \bmod p \\
 k_{AW} = k_1 = k_{WA}
 \end{array}
 \quad
 \begin{array}{l}
 k_{WA} = (t_A)^j \bmod p \\
 = (g^u)^j \bmod p \\
 = g^{uj} \bmod p
 \end{array}
 \quad
 \begin{array}{l}
 k_{WB} = (t_B)^i \bmod p \\
 = (g^v)^i \bmod p \\
 = g^{vi} \bmod p \\
 k_{WB} = k_2 = k_{BW}
 \end{array}
 \quad
 \begin{array}{l}
 k_{BW} = (W_A)^v \bmod p \\
 = (g^i)^v \bmod p \\
 = g^{iv} \bmod p
 \end{array}$$

### Encryption based KAP - EKAP



$k = \text{randi}(Z_p^*)$   
 $c_k = \text{Enc}(\text{PuK}_B, i, k)$   
**M**-message to be encrypted with symmetric encryption method e.g. AES128  
 $C_k = E(k, M) = \text{AES128}(k, e, M)$

$\xrightarrow{C_k}$

$k = \text{Dec}(\text{PrK}_B, C_k)$   
 $M = D(k, d, C_k) = \text{AES128}(k, d, C_k)$

### MiMA



Imagine that *W* generated her  $\text{PrK}_W = z$  and  $\text{PuK}_W = e$ .  
*W* send a message to *A* writing the following message:  
 "Dear *A* I am *B* and I am sending you my  $\text{PuK} = e$  for our further communications. Trully yours *B*."

<http://crypto.fmf.ktu.lt/xdownload/>

- [octave-7.3.0-w64-installer.exe](#)
- [octave.Stud.7z](#)

```

GNU Octave, version 7.1.0
Copyright (C) 1993-2022 The Octave Project
This is free software; see the source code
There is ABSOLUTELY NO WARRANTY; not even
FITNESS FOR A PARTICULAR PURPOSE. For details
see the GNU General Public License at
https://www.gnu.org/licenses/gpl.html

Octave was configured for "x86_64-w64-mingw64".

Additional information about Octave is available at
https://www.octave.org.

Please contribute if you find this software useful.
For more information, visit https://www.octave.org.

Read https://www.octave.org/bugs.html to
report for information about changes from previous
versions.

>> 2^8
ans = 256
>> e = mod_exp(2,8,10)
e = 6
>> p = genstrongprime(28)
p = 232702259
>> pb = dec2bin(p)
pb = 1101110111101100000100110011
>>

```

```

AES128.m
19 % Encryption example:
20 >> in = 'Hello Bob';
21 >> kh32 = '000000000000000000000000099828F0';
22 >> NR = 10;
23 >> Ch = AES128(in, kh32, NR, 'e');
24 ASCII_e = 57:1:57-mV % ciphertext in ASCII format
25 Ch = 0f9a2c08d191310fb27ed16d90f45686 % ciphertext in hexadecimal format
26 %
27 % Decryption example:
28 >> Dh = AES128(Ch, kh32, NR, 'd');
29 Dh = 00000000000004865c6c6f7720426f62 % decrypted message in hex format
30 D = Hello Bob % Decrypted message in ASCII format
31 %
32 function Out = AES128(in, key, Nr, mode)

```

Private and Public keys generation :  $PrK = x$  ;  $PuK = a$  ;

1) Generate strong prime number  $P$ .

$\gg p = \text{genstrongprime}(28)$  % generates 28 bit lengths of  $p$

2) Find a generator  $g$  in the set  $\mathbb{Z}_p^* = \{1, 2, 3, \dots, p-1\}$

If  $p$  is strong prime, then  $p = 2 \cdot q + 1$ , when  $q$  is also prime:  $11 = 2 \cdot 5 + 1$ ;  $23 = 2 \cdot 11 + 1$ .

$\gg q = (p-1)/2$

$\gg g = 2$

$\gg \text{mod\_exp}(g, q, p) \neq 1$  % I-st condition  
 % If it is equal to 1  $\rightarrow$  choose the other  $g$   
 % If no, then verify;

$\gg \text{mod\_exp}(g, 2, p) \neq 1$  % II-nd condition  
 % If it is equal to 1  $\rightarrow$  choose the other  $g$ .

3) Generate  $PrK = x$  using random number generator function  $\text{randi}$

$\gg x = \text{int64}(\text{randi}(2^{28}-1))$   $\gg x = \text{randi}(2^{28}-1)$

3) generate  $PrK = x$  using random number generator function randi

$\Rightarrow x = \text{int64}(\text{randi}(2^{28}-1))$

```
>> x=randi(2^28-1)
```

```
x = 1.9906e+08
```

```
>> x=int64(randi(2^28-1))
```

```
x = 256210849
```

4) compute  $PuK = a$  using DEF, i.e. function

$\Rightarrow a = \text{mod\_exp}(g, x, p)$

$PrK_A = x \leftarrow \text{randi} \Rightarrow PuK_A = a = g^x \bmod p$

$PrK_B = y \leftarrow \text{randi} \Rightarrow PuK_B = b = g^y \bmod p$

For example:

```
>> k=randi(2^28-1)
```

```
k = 3.6754e+07
```

```
>> k=int64(randi(2^28-1))
```

```
k = 233360567
```

```
>> kb=dec2bin(k)
```

```
kb = 1101111010001100110010110111
```

```
>> maxb=dec2bin(2^28-1)
```

```
maxb = 1111111111111111111111111111
```

```
>> max=int64(2^28-1)
```

```
max = 268435455
```

```
>> p=genstrongprime(28)
```

```
p = 227317067
```

```
>> isprime(p)
```

```
ans = 1
```

```
>> q=(p-1)/2
```

```
q = 113658533
```

```
>> isprime(q)
```

```
ans = 1
```

```
>>
```

```
>> g=2
```

```
g = 2
```

```
>> mod_exp(g,q,p)
```

```
ans = 227317066
```

```
>> mod_exp(g,2,p)
```

```
ans = 4
```

```
>> x=int64(randi(2^28-1))
```

```
x = 125675623
```

```
>> a=mod_exp(g,x,p)
```

```
a = 67591766
```